

# The Key to Effective UDF Optimization: Before Inlining, First Perform Outlining

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**Carnegie Mellon**  
Database Group

## VLDB 2025 Runner-Up Best Paper Award

# UDFs are Popular!



```
SELECT customer_name, is_vip(customer_key)
FROM customer;
```


SQL

# UDFs are Popular!



```
SELECT customer_name, is_vip(customer_key)
FROM customer;
```

SQL



```
CREATE FUNCTION is_vip(ckey INTEGER)
RETURNS BOOLEAN DECLARE DECIMAL total; BOOLEAN vip;
total = (SELECT SUM(o_totalprice) FROM orders
        WHERE o_custkey = ckey);
IF (total > 1000000) THEN vip = True;
    ELSE vip = False;
RETURN vip;
```

UDF

# UDFs are Popular!



```
SELECT customer_name, is_vip(customer_key)
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SQL

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CREATE FUNCTION is_vip(ckey INTEGER)
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RETURN vip;
```

UDF

[VLDB 2018]

**Billions of queries  
per day invoke  
UDFs**



# UDF Optimization is Hard!



```
SELECT customer_name, is_vip(customer_key)
FROM customer;
```

SQL



UDF

???

**UDFs are opaque  
to the query  
optimizer!**

# UDF Inlining (FROID)



**UDF**

**Hard to optimize**



# UDF Inlining (FROID)



**UDF**

**Hard to optimize**



**VLDB 2018**

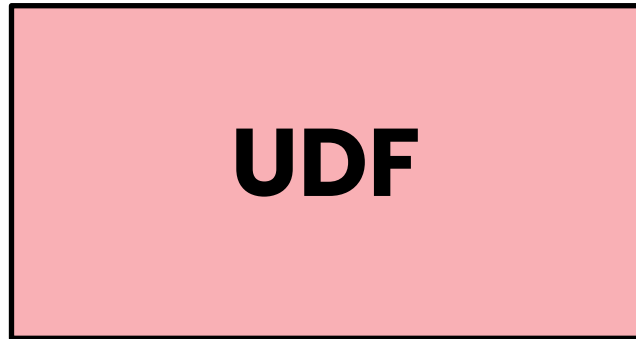
**Let's translate  
UDFs to SQL**

**Froid: Optimization of  
Imperative Programs in  
a Relational Database**



**Microsoft**

# UDF Inlining (FROID)



**UDF Inlining**



**Hard to optimize**



**VLDB 2018**

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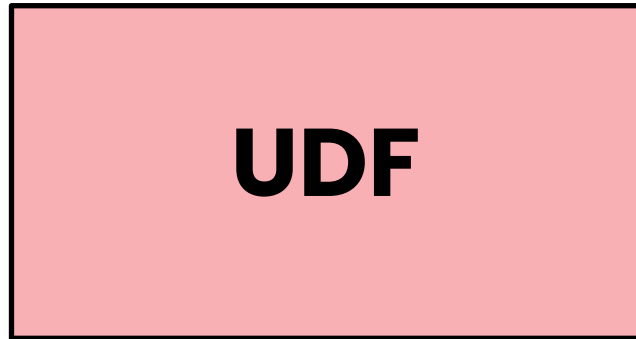
**Froid: Optimization of  
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**Microsoft**



# UDF Inlining (FROID)



**UDF Inlining**



**Hard to optimize**



**Easy to optimize**



**Let's translate  
UDFs to SQL**

**VLDB 2018**

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**Microsoft**

# Talk Overview



# Talk Overview



Inlining leads to **sub-optimal performance**

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Inlining leads to **sub-optimal performance**

Inlining **entire** UDFs creates **complex** queries

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Instead, we inline only the important **pieces** of a UDF

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We achieve this through **UDF outlining**

# Talk Overview



Inlining leads to **sub-optimal performance**

Inlining **entire** UDFs creates **complex** queries

Instead, we inline only the important **pieces** of a UDF

We achieve this through **UDF outlining**

Our approach outperforms inlining by more than **1000x**

# UDF Inlining (FROID)



```
total = (SELECT ...);
```

UDF

```
IF (total > 1000000) THEN vip = True;
```

```
    ELSE vip = False;
```

```
RETURN vip;
```

**UDF Inlining**

A large red arrow pointing to the right, indicating the direction of the UDF inlining process.



# UDF Inlining (FROID)



```
total = (SELECT ...);
```

UDF

```
IF (total > 1000000) THEN vip = True;
```

```
    ELSE vip = False;
```

```
RETURN vip;
```

**UDF Inlining**

**Chain together translated UDF statements with LATERAL joins**

# UDF Inlining (FROID)



```
total = (SELECT ...);
```

UDF

```
IF (total > 1000000) THEN vip = True;  
    ELSE vip = False;
```

```
RETURN vip;
```

UDF Inlining

```
SELECT T2.vip FROM
```

SQL

```
(SELECT (SELECT ...) AS total) T1
```

LATERAL

```
(SELECT CASE WHEN (T1.total > 1000000)  
    THEN True ELSE False  
    END AS vip) T2
```

Chain together translated UDF  
statements with LATERAL joins

# The Problem with Inlining



```
SELECT customer_name,
```

SQL

```
SELECT T2.vip FROM
```

SQL

```
(SELECT (SELECT ...) AS total) T1
```

LATERAL

```
(SELECT CASE WHEN (T1.total > 1000000)
      THEN True ELSE False
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```

```
FROM customer;
```

# The Problem with Inlining



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LATERAL

```
(SELECT CASE WHEN (T1.total > 1000000)
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```

```
FROM customer;
```

**Complex subqueries!**

# The Problem with Inlining



```
SELECT customer_name,
```

SQL

```
SELECT T2.vip FROM
```

SQL

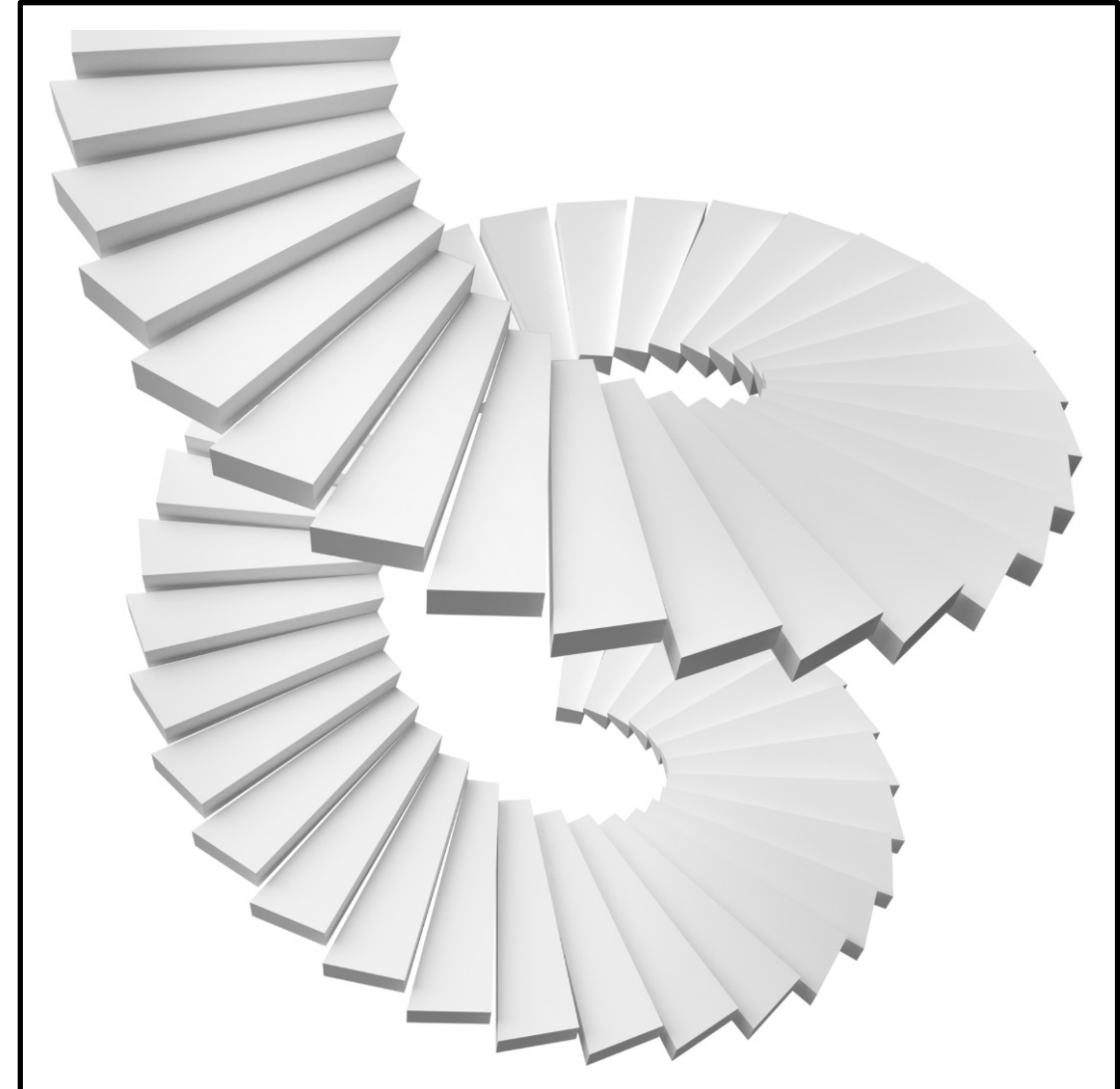
```
(SELECT (SELECT ...) AS total) T1
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LATERAL

```
(SELECT CASE WHEN (T1.total > 1000000)
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```
FROM customer;
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**Complex subqueries!**



# The Problem with Inlining



```
SELECT customer_name,
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SELECT T2.vip FROM
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SQL

```
(SELECT (SELECT ...) AS total) T1
```

LATERAL

```
(SELECT CASE WHEN (T1.total > 1000000)
      THEN True ELSE False
      END AS vip) T2
```

```
FROM customer;
```

```
SELECT customer_name, ...
```

SQL

```
FROM customer
```

```
LEFT OUTER JOIN orders ...
```

## Subquery Unnesting



# The Problem with Inlining



```
SELECT customer_name,
```

SQL

```
SELECT T2.vip FROM
```

SQL

```
(SELECT (SELECT ...) AS total) T1
```

LATERAL

```
(SELECT CASE WHEN (T1.total > 1000000)
      THEN True ELSE False
      END AS vip) T2
```

```
FROM customer;
```

```
SELECT customer_name, ...
FROM customer
LEFT OUTER JOIN orders ...
```

SQL

**Subquery Unnesting**



**Inlining generates complex subqueries  
that database systems can't unnest!**

# The Problem with Inlining



		ProcBench Queries														
Technique		Q1	Q5	Q6	Q7	Q9a	Q9b	Q12	Q13	Q15	Q17	Q18	Q20a1	Q20a2	Q20b1	Q20b2
SQL Server	Inlining	×	×	×	×	✓	✓	×	×	×	×	×	✓	✓	×	×



**Inlining: 4/15**

**Inlining generates complex subqueries  
that database systems can't unnest!**



# Why Inline the Entire UDF?



```
total = (SELECT ...);
```

UDF

```
IF (total > 10000000) THEN vip = True;  
                               ELSE vip = False;  
  
RETURN vip;
```

# Why Inline the Entire UDF?



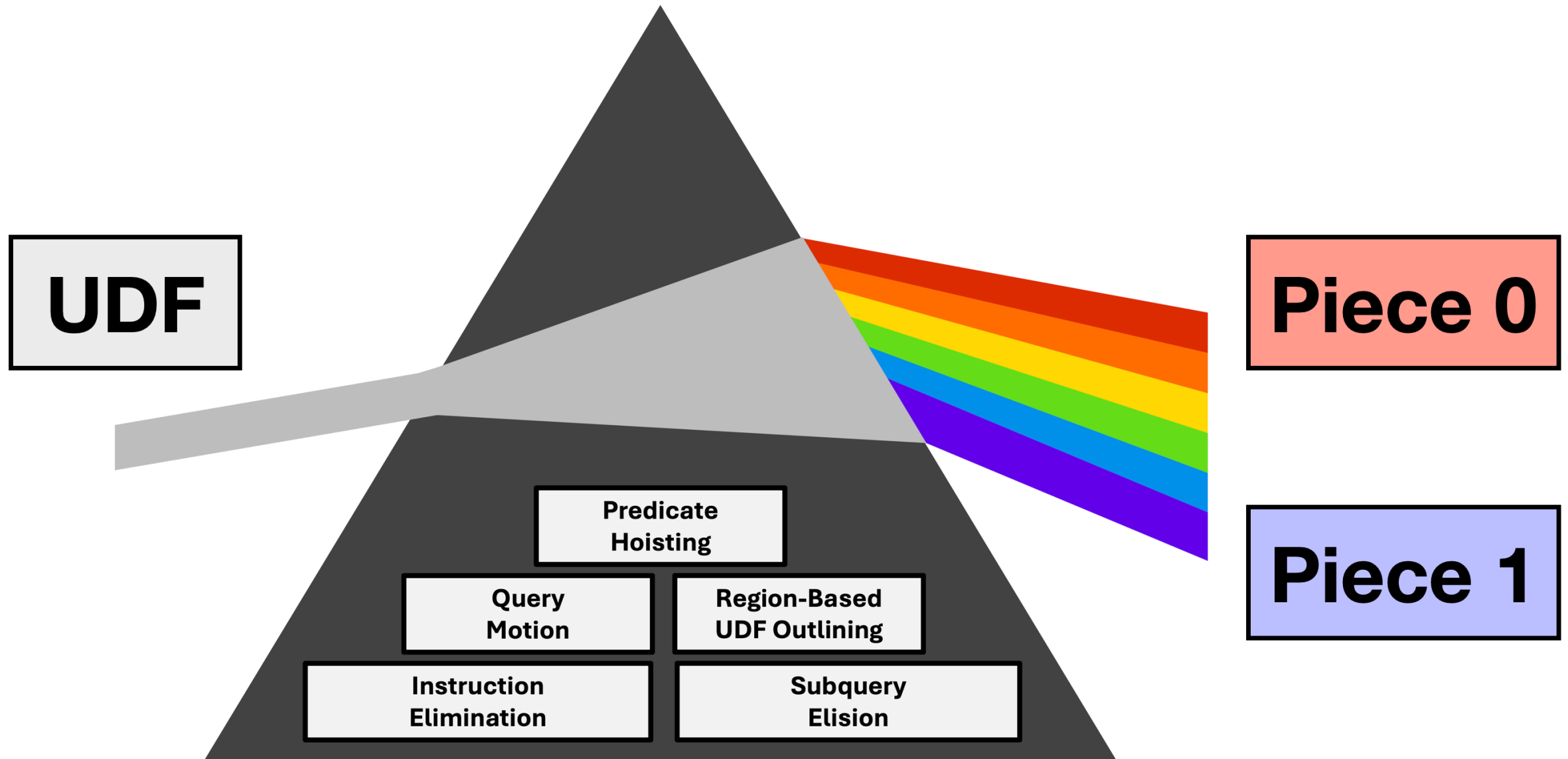
```
total = (SELECT ...);
```

UDF

```
IF (total > 10000000) THEN vip = True;  
                               ELSE vip = False;  
  
RETURN vip;
```

**Can we "outline" irrelevant UDF code?**

# PRISM: A UDF Optimizing Compiler



# PRISM: A UDF Optimizing Compiler

**P**redicate Hoisting

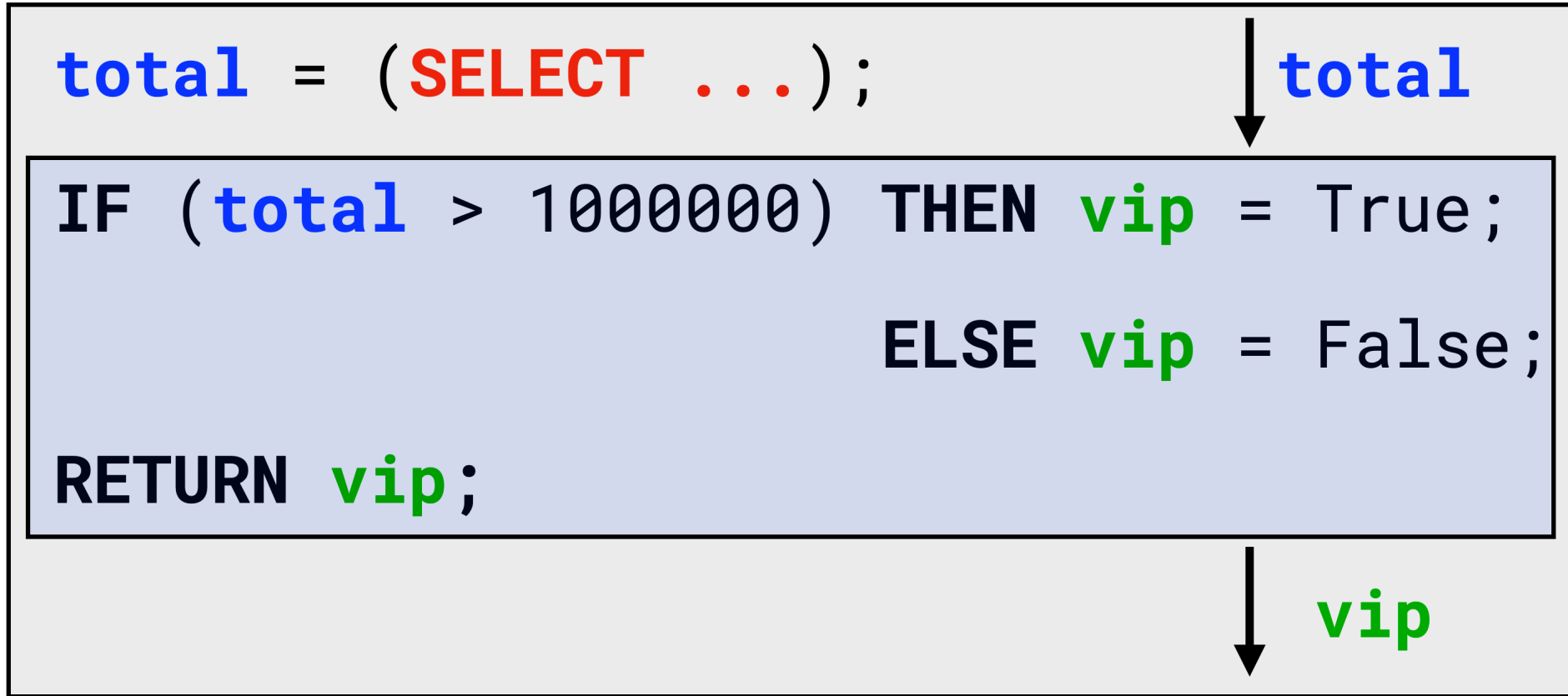
**R**egion-Based UDF Outlining

**I**nstruction Elimination

**S**ubquery Elision

Query **M**otion

# Region-Based UDF Outlining



**Identify regions of irrelevant UDF code**

# Region-Based UDF Outlining



```
total = (SELECT ...);
```

total



```
IF (total > 1000000) THEN vip = True;
```

```
ELSE vip = False;
```

```
RETURN vip;
```

```
RETURN f(total);
```

vip



**Outline regions into separate functions**

# Region-Based UDF Outlining



```
total = (SELECT ...);
```

```
RETURN f(total);
```

**Replace regions with function calls**

# Instruction Elimination



```
total = (SELECT ...);  
RETURN f(total);
```

**Eliminate redundant instructions**



# Instruction Elimination



```
RETURN f(SELECT ...);
```

**Eliminate redundant instructions**

# Instruction Elimination



```
RETURN f(SELECT ...);
```

**Eliminate redundant instructions**

# Subquery Elision



```
RETURN f(SELECT ...);
```



```
SELECT customer_name, is_vip(customer_key)  
FROM customer;
```

SQL

**Directly inject the UDF's return value**

# Subquery Elision



```
RETURN f(SELECT ...);
```



```
SELECT customer_name, f(SELECT ...)  
FROM customer;
```

SQL

**Directly inject the UDF's return value**

# Subquery Elision



```
SELECT customer_name, f(SELECT ...)  
FROM customer;
```

SQL

**Directly inject the UDF's return value**

# Inlining vs PRISM



```
SELECT customer_name,
```

SQL

```
SELECT T2.vip FROM
```

SQL

```
(SELECT (SELECT ...) AS total) T1
```

LATERAL

```
(SELECT CASE WHEN (T1.total > 1000000)
      THEN True ELSE False
      END AS vip) T2
```

```
FROM customer;
```

```
SELECT customer_name, f(SELECT ...)
```

```
FROM customer;
```

SQL

**PRISM generates simpler, faster queries**

# Evaluation



**Ran PRISM-optimized UDFs in DuckDB & SQL Server**

# Evaluation



**Ran PRISM-optimized UDFs in DuckDB & SQL Server**

**Ran the ProcBench @ 10GB Scale Factor (TPC-DS + UDFs)**

**VLDB 2021**



**Procedural Extensions of SQL:  
Understanding Their Usage in the Wild**



# Evaluation



**Ran PRISM-optimized UDFs in DuckDB & SQL Server**

**Ran the ProcBench @ 10GB Scale Factor (TPC-DS + UDFs)**

**Intel Xeon Gold 5218R CPU, 192GB DDR4 RAM, 960GB NVMe SSD**

**VLDB 2021**



**Microsoft**

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**Built Column-Store Indexes on SQL Server**

**VLDB 2021**



**Microsoft**

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**Ran PRISM-optimized UDFs in DuckDB & SQL Server**

**Ran the ProcBench @ 10GB Scale Factor (TPC-DS + UDFs)**

**Intel Xeon Gold 5218R CPU, 192GB DDR4 RAM, 960GB NVMe SSD**

**Built Column-Store Indexes on SQL Server**

**2x Cold Runs and 5x Hot Runs**

**VLDB 2021**



**Microsoft**

**Procedural Extensions of SQL:  
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# Evaluation (Unnesting)



		ProcBench Queries														
Technique		Q1	Q5	Q6	Q7	Q9a	Q9b	Q12	Q13	Q15	Q17	Q18	Q20a1	Q20a2	Q20b1	Q20b2
SQL Server	Inlining	×	×	×	×	✓	✓	×	×	×	×	×	✓	✓	×	×

Inlining: 4/15

# Evaluation (Unnesting)



Technique		ProcBench Queries														
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SQL Server	Inlining	✗	✗	✗	✗	✓	✓	✗	✗	✗	✗	✗	✓	✓	✗	✗
	PRISM	✓	✓	✗	✓	✓	✓	✓	✓	✗	✓	✗	✓	✓	✓	✓



**Inlining: 4/15**

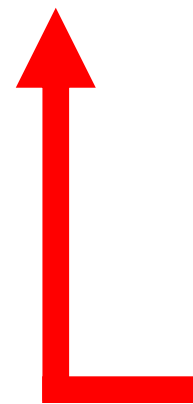
**PRISM: 12/15**

# Evaluation (Unnesting)



		ProcBench Queries														
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SQL Server	Inlining	✗	✗	✗	✗	✓	✓	✗	✗	✗	✗	✗	✓	✓	✗	✗
	PRISM	✓	✓	✗	✓	✓	✓	✓	✓	✗	✓	✗	✓	✓	✓	✓
DuckDB	Inlining	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓
	PRISM	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓

**BTW 2015**  
**Unnesting  
Arbitrary Queries**



**Inlining: 4/15**

**PRISM: 12/15**

**DuckDB: 15/15**

# Evaluation (Unnesting)



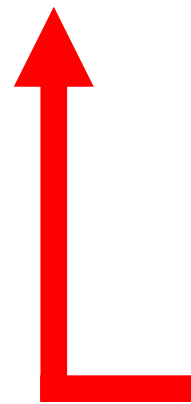
Add support for nested laterals #7528

Merged Mytherin merged 4 commits into `duckdb:feature` from `CMU-15-745:nested_laterals`

DuckDB	Inlining PRISM	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓
		✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓

BTW 2015

Unnesting  
Arbitrary Queries



Inlining: 4/15

PRISM: 12/15

DuckDB: 15/15



# Evaluation (Overall Speedup)



	Average Speedup	Maximum Speedup
<b>SQL Server</b>	298.7×	
<b>DuckDB</b>		



**Unnesting**



# Evaluation (Overall Speedup)



	<b>Average Speedup</b>	<b>Maximum Speedup</b>
<b>SQL Server</b>	298.7×	2997.9×
<b>DuckDB</b>		



**Unnesting**

# Evaluation (Overall Speedup)



	<b>Average Speedup</b>	<b>Maximum Speedup</b>
<b>SQL Server</b>	298.7×	2997.9×
<b>DuckDB</b>	1.3×	



**Fewer LATERAL Joins**

# Evaluation (Overall Speedup)



	<b>Average Speedup</b>	<b>Maximum Speedup</b>
<b>SQL Server</b>	298.7×	2997.9×
<b>DuckDB</b>	1.3×	2270.2×



**Better Query Plan**

# Evaluation (Overall Speedup)



	Average Speedup	Maximum Speedup
<b>SQL Server</b>	298.7×	2997.9×
<b>DuckDB</b>	1.3×	2270.2×



Correctly visit all expressions during lateral join decorrelation, particularly with nested lateral joins #10936

Merged Mytherin merged 2 commits into `duckdb:main` from `Mytherin:correlatedsubquerybinding`

# Evaluation (Overall Speedup)



	<b>Average Speedup</b>	<b>Maximum Speedup</b>
<b>SQL Server</b>	298.7×	2997.9×
<b>DuckDB</b>	1.3×	18.0×



**Compiled Loop vs Recursive SQL**

# Conclusion



# Conclusion



Inlining leads to **sub-optimal performance**

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Inlining **entire** UDFs creates **complex** queries



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We achieve this through **UDF outlining** with **PRISM**

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Inlining leads to **sub-optimal performance**

Inlining **entire** UDFs creates **complex** queries

Instead, we inline only the important **pieces** of a UDF

We achieve this through **UDF outlining** with **PRISM**

Our approach outperforms inlining by more than **1000x**

# Future Work



**Python UDFs.**

samarch.xyz